



Discrete Mathematics

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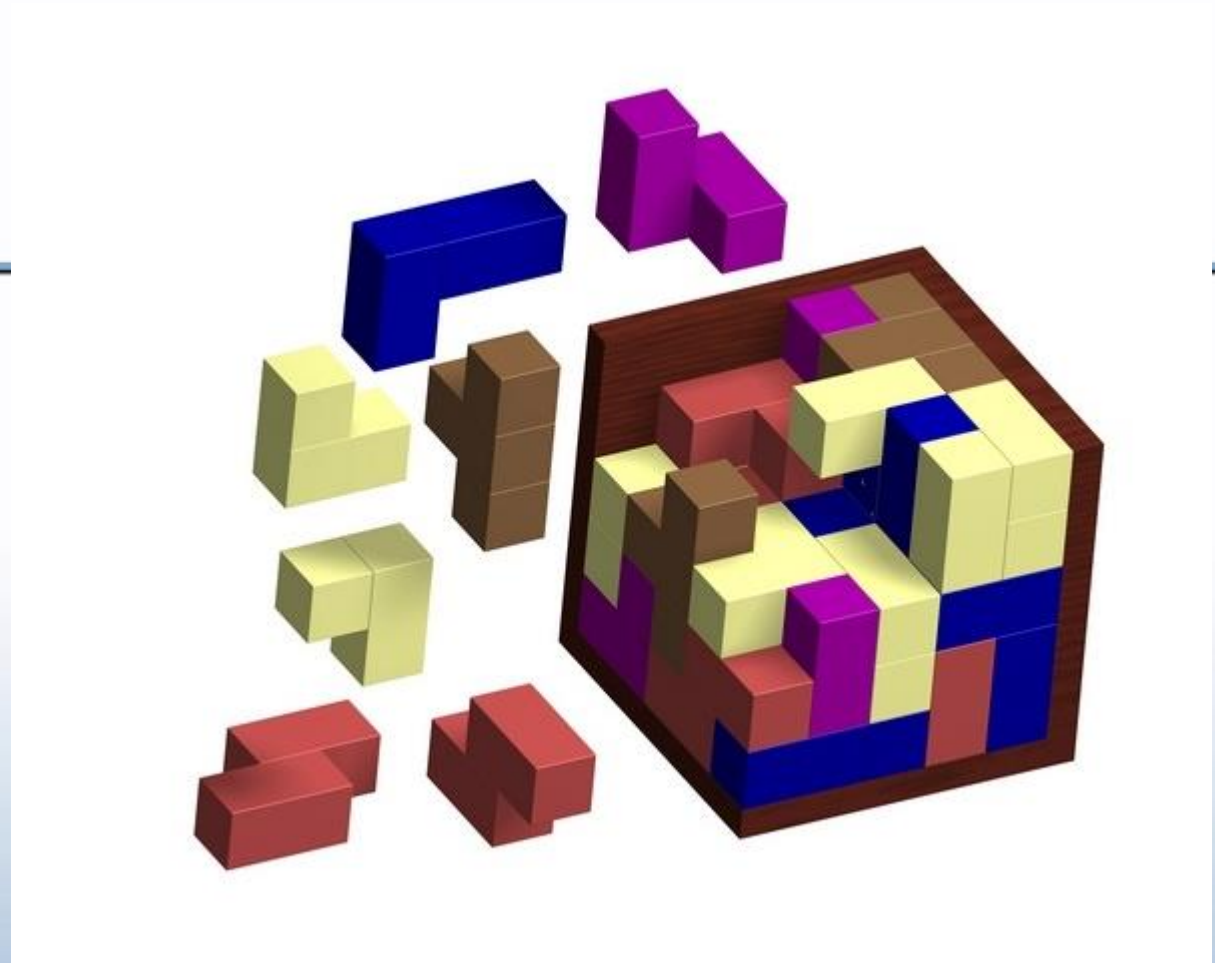
Unit materials

- Lecture notes
- Seminar handouts

are available at

<http://gm.softalliance.net/>

Advice: download and print lecture notes
before the next lecture



Propositional Logic

<http://www.craftsmanspace.com>



Contents

- Notions of logic
- Branches of logic
- Propositional logic
- Boolean operators
- Logic and computing
- Equivalence rules and derivations



Notion of Logic

The word *logic* is inherited from classical Greek λόγος logos; meaning word, thought, idea, argument, account, reason, or principle.

In common sense, logic is a tool for distinguishing between the true and the false (Averroes).

Logic is the study of the principles and criteria of valid inference and demonstration.



Inference means making a conclusion based solely on what one already knows.

Inference is the process of deriving logical conclusions from premises known or assumed to be true.

Example

- | | |
|---------------------------|--------------|
| Flipper is a dolphin | - premise |
| Every dolphin is a mammal | - premise |
| Flipper is a mammal | - conclusion |



Branches of Logic

- **Informal logic** is the study of natural language arguments.
- **Argument** is a set of sentences known as the premises, and another sentence known as the conclusion in which it is assumed that the truth of the conclusion follows from the premises.

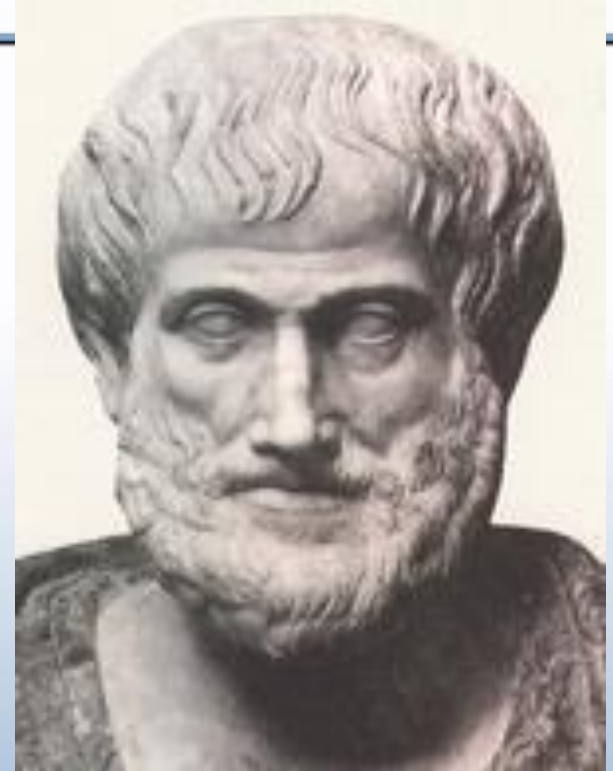


- **Formal logic** is the study of inference with purely abstract content not related to any particular thing or property.
- **Mathematical logic** is the study of formal features of logical inference: mathematical study of logic and the application of this study to other areas of mathematics.



Formal Logic

- **Subject:** definition of rules of proper inference (when correctly applied to true premises, lead to true conclusions).
- The first rules of formal logic were written by **Aristotle** in his **Organon**.



Aristotle
(384-322 BC)



Syllogisms

- Aristotle defined a number of **syllogisms**, correct three-part inferences, that can be used as building blocks for more complex reasoning.

Example: most famous syllogism

All men are mortal. Socrates is a man.

Therefore Socrates is mortal.

Two premises and conclusion here are true.



- Structure of a syllogism
 - Major premise - a general statement
 - Minor premise - a specific statement
 - Conclusion - based on the two premises

Each part is a proposition in the form:

"All A are B," "Some A are B",
"No A are B" or "Some A are not B"

There are 24 types of logically distinct **valid** syllogisms.



Does the truth of the conclusion follow from truth of the premises?

Example: valid inference with false premises

- All apples are blue. (False)
- A banana is an apple. (False)
- Therefore, a banana is blue. (False)

For the conclusion to be necessarily true, the premises need to be true.



Example: invalid inference

- All A are B.
- C is a B.
- Therefore, C is an A.

This is invalid form of inference, because from true premises it can lead to a false conclusion:

- All apples are fruit. (True)
- Bananas are fruit. (True)
- Therefore, bananas are apples. (False)



Properties of inference:

- For the conclusion to be necessarily true, the premises need to be true.
- An inference can be valid even if the parts are false, and can be invalid even if the parts are true.
- A valid inference does not depend on the truth of the premises and conclusion, but on the rules of inference studied in formal logic.
- A valid form of inference with true premises will always have a true conclusion.

Mathematical Logic

- Logic is the basis for all mathematical reasoning. To be able to understand and construct our own correct mathematical arguments we must understand logic.
- **Mathematical Logic** studies formal features of logical inference using symbolic abstractions.
- Mathematical Logic is
 - a tool for working with compound statements built from simpler statements.
 - the foundation for expressing formal proofs in all branches of mathematics.



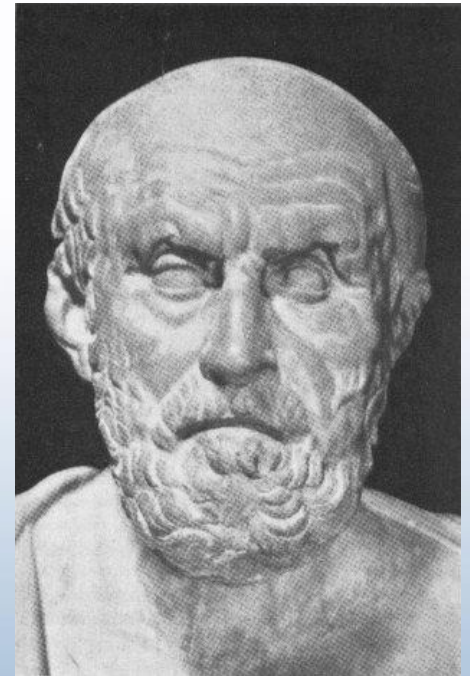
- Mathematical Logic includes:
 - a formal language for expressing compound statements
 - a concise notation for writing them
 - a methodology for objectively reasoning about their truth or falsity.
- Divided into two branches:
propositional logic and *predicate logic*.



Propositional Logic

Aristotle developed a detailed system of logic and Chrysippus of Soli introduced **propositional logic** centered around logical operations.

It is the logic of compound statements built from simpler statements using **logic operators** (NOT, AND, OR).



Chrysippus of Soli
(ca. 281 B.C. – 205 B.C.)



Some applications in computer science:

- Design of digital electronic circuits
- Construction of computer programs
- Verification of the correctness of programs
- Queries to databases & search engines
- **Constructive geometric modeling**



Definition of a Proposition

- A **proposition** is:
 - ❖ a *statement* (a declarative sentence)
 - with some definite meaning (not vague or ambiguous)
 - ❖ having a **truth value** that is either **true** or **false** (under interpretation)
 - it is never both, neither, or somewhere “in between”
 - however, you might not know the actual truth value and the truth value might depend on the situation or context.



Propositions in Natural Language

Propositions:

- “It is raining.”
- “London is the capital of China.”
- $1 + 2 = 2$

NOT propositions:

- “Who is there?” (interrogative: no truth value)
- “ $1 + 2$ ” (term: no truth value)
- “kudliva bokra” (no definite meaning in known languages)
- $7 - z = 77$ (neither true nor false, since variables are not assigned with their values)
- “Go to the town!” (imperative: no truth value)

Propositions in Natural Language



Q: Determine whether the following statements are propositions or not and explain your answers:

- a) The Earth is flat.
- b) Is it raining?
- c) Stop and give way to pedestrians.
- e) $2 + 8 = 10$
- d) $2x - 7 = 8$

Formally not a proposition because it is neither true nor false. Note that it can be turned into a proposition if we assign a value to the variable x .



Propositional Variables

- Propositional variables (statement variables, atoms) are simply letters which represent propositions: $p, q, r, s, t \dots$
Correspond to simple (English) sentences.
Examples:
Proposition p is “I had salad for lunch”
Proposition q is “Today is Monday”
❖ Proposition s is
❖ Proposition t is ...



Truth Value

- Proposition is either true or false but not both
- Each proposition has a **Truth Value**.
- If a proposition is
 - true, we denote that by **T**
 - false, we denote that by **F**

T \equiv **True**; **F** \equiv **False**

“ \equiv ” means “ is defined as ”

❖ Proposition p is “Sun is a planet”. It is



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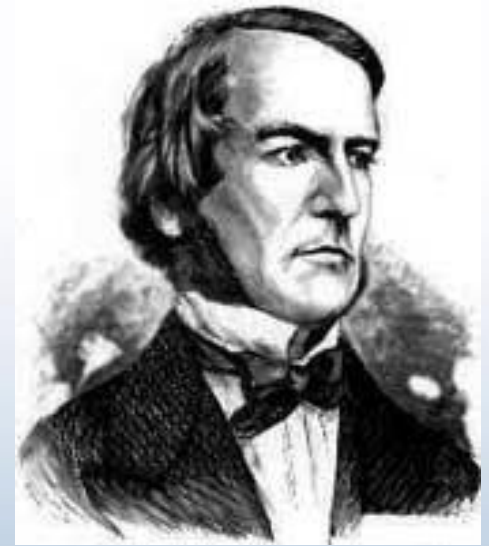
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Compound Propositions

Compound statements are built from simpler statements using **logic operators** or **Boolean connectives**.

These methods were discussed by the English mathematician George Boole in 1854 in his book “The Laws of Thought”.



George Boole
(1815-1864)

Compound Propositions



- To create new propositions we may combine one, two or more propositions into **complex (compound) propositions**.
- Compound propositions are built up from atoms using operators such as NOT, AND, OR. Correspond to compound English sentences (“I had salad for lunch **AND** I had a steak for dinner.”)



Operators / Connectives

- **Operator or connective** combines n operands (expressions) into a larger expression, *e.g.*, “+” in numeric expression.
- **Unary operators** take 1 operand (*e.g.*, -3)
Binary operators take 2 operands (*e.g.*, 3×4)
- **Propositional or Boolean operators (connectives)** operate on propositions instead of numbers.



Some Boolean Operators

<u>Formal Name</u>	<u>Name</u>	<u>Arity</u>	<u>Symbol</u>
Negation operator	NOT	Unary	\neg
Conjunction operator	AND	Binary	\wedge
Disjunction operator	OR	Binary	\vee
Exclusive-OR operator	XOR	Binary	\oplus



Negation Operator

- The unary **negation operator** “ \neg ” (logical *NOT*) transforms a proposition into its **negation**.
- The proposition $\neg p$ is read “not p.”
- The truth value of $\neg p$, is the opposite of the truth value of p .

Example: If p = “I have brown hair.”

then $\neg p$ = “I do **not** have brown hair.”



Truth Table for Negation Operator

- A **truth table** displays the relationships between the truth values of propositions, when applying an operator to them.

Truth table for proposition and it's negation:

Operand
column

Result
column

p	$\neg p$
T	F
F	T



Negation Exercise

Find the negation of the proposition

p = "At least 10 inches of rain fell today in Poole."

and express this in simple English.

Solution: the negation is

$\neg p$ = "It is not the case that at least 10 inches of rain fell today in Poole."

This negation can be more simply expressed by

$\neg p$ = "Less than 10 inches of rain fell today in Poole."



Conjunction Operator

Binary conjunction operator “ \wedge ” (logical *AND*) combines two propositions to form their logical conjunction.

The conjunction $p \wedge q$ is true when both p and q are true and is false otherwise.

Example:

If p = “I had salad for lunch.” and

q = “I had a steak for dinner.”,

then $p \wedge q$ = “I had salad for lunch
and I had a steak for dinner.”



Conjunction Truth Table

- Note:
a conjunction
 $p_1 \wedge p_2 \wedge \dots \wedge p_n$
of n propositions
will have 2^n rows
in its truth table.

Operand Columns		Result Column
p	q	$p \wedge q$
F	F	F
F	T	F
T	F	F
T	T	T

- Note: \neg **NOT** and \wedge **AND** operations together are sufficient to express any Boolean truth table!



Conjunction Exercise

For the two given propositions

p ="Today is Tuesday."

q ="It is raining today."

when the conjunction is false?

Solution: the conjunction

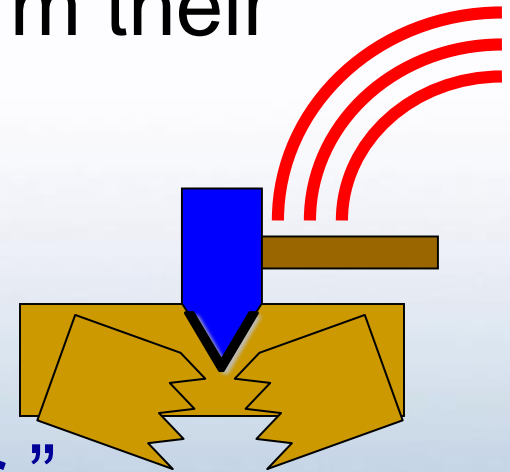
$p \wedge q$ = "Today is Tuesday and
it is raining today."

is false on any day that is not a Tuesday and on
Tuesdays when it does not rain.



Disjunction Operator

- Binary disjunction operator “ \vee ” (logical *OR*) combines two propositions to form their logical disjunction.
- Example:
- p = “My car has a bad engine.”
- q = “My car has a bad carburettor.”
- $p \vee q$ = “Either my car has a bad engine, **or** my car has a bad carburettor.”



After the downward-pointing “axe” of “ \vee ” splits the wood, you can take 1 piece **OR** the other, or both.



Disjunction Truth Table

- Note that $p \vee q$ means that p is true, or q is true, **or both** are true!
- This operation is also called **inclusive or**, because it **includes** the possibility that both p and q are true.
- Note: “ \neg ” NOT and “ \vee ” OR together are also universal.

p	q	$p \vee q$
F	F	F
F	T	T
T	F	T
T	T	T

Note difference from AND



Disjunction Exercise

For the two given propositions

p ="Today is Tuesday."

q ="It is raining today."

when the disjunction is false?

Solution: the disjunction

$p \vee q$ = "Today is Tuesday or it is raining today."

is only false on days that are not Tuesdays when it also does not rain.



Nested Propositional Expressions

- Use parentheses to **group sub-expressions**:
“I just saw my old friend, and either he’s grown or I’ve shrunk.” $:\equiv f \wedge (g \vee s)$

$(f \wedge g) \vee s$ would mean something different

$f \wedge g \vee s$ would be ambiguous

- By convention, “ \neg ” NOT takes **precedence** over both “ \wedge ” and “ \vee ”.

$\neg s \wedge f$ means $(\neg s) \wedge f$, not $\neg (s \wedge f)$



Simple Exercise

p = “It rained last night”,

q = “The sprinklers came on last night,”

r = “The lawn was wet this morning.”

Translate each of the following into English:

- $\neg p$ = “It didn’t rain last night.”
- $r \wedge \neg p$ = “The lawn was wet this morning, and it didn’t rain last night.”
- ❖ $r \wedge (p \vee q)$ =



Exclusive-Or Operator

- Binary **exclusive-or operator** “ \oplus ” (logical *XOR*) combines two propositions to form their logical “exclusive or” (exjunction).

Example:

- p = “I will earn an A in this course,”
- q = “I will drop this course,”
- $p \oplus q$ = “I will either earn an A in this course, or I will drop it (but not both!) ”



Exclusive-Or Truth Table

- Note that $p \oplus q$ means that p is true, or q is true, but **not both!**

- This operation is called **exclusive or**, because it **excludes** the possibility that both p and q are true.

p	q	$p \oplus q$
F	F	F
F	T	T
T	F	T
T	T	F

Note
difference
from OR.



Exclusive-Or Exercise

For the two given propositions

p ="Today is Tuesday."

q ="It is raining today."

when the exclusive-or is false?

Solution: the exclusive-or

$p \oplus q$ = "Today is either Tuesday or it is raining today,
but not both."

is false on rainy Tuesdays and on any other weekday,
when it does not rain.







Boolean Operators Summary

p	q	$\neg p$	$p \wedge q$	$p \vee q$	$p \oplus q$
F	F	T	F	F	F
F	T	T	F	T	T
T	F	F	F	T	T
T	T	F	T	T	F



Some Alternative Notations

Name:	not	and	or	xor
Propositional logic:	\neg	\wedge	\vee	\oplus
Boolean algebra:	\bar{p}	pq	$+$	\oplus
C/C++/Java (wordwise):	!	& &		!=
C/C++/Java (bitwise):	~	&		^
Logic gates:				



Implication

- Implication (conditional statement) $p \rightarrow q$ is false when p is true and q is false, and true otherwise.
- p is **hypothesis** (or **antecedent** or **premise**) and q is **conclusion** (or **consequence**).
- Note: The implication is false only when P is true and Q is false!



- Equivalent forms:

- If p , then q
- p implies q
- If p , q
- p only if q
- p is a sufficient condition for q
- q if p
- q whenever p
- q is a necessary condition for p

p	q	$p \rightarrow q$
F	F	T
F	T	T
T	F	F
T	T	T



Implication Example

Professor's promise:

$p \rightarrow q$: “If you get 100% on the final, then you will get an A.”

- If you manage to get a 100% on the final, then you would expect to receive an A: **promise is kept**
- If you do not get 100% you may or may not receive an A depending on other factors: **promise can be kept**
- However, if you do get 100%, but the professor does not give you an A, you will feel cheated: **false promise**



Conditional Statement in Programming

Let a program statement take FALSE value, if there is an error (division by zero, etc.).

Conditional statement $A \rightarrow B$ in the form

if(A) then B;

triggers an error event (FALSE value), only if

the condition A is TRUE and

B is FALSE,

meaning an error happens when executing B.



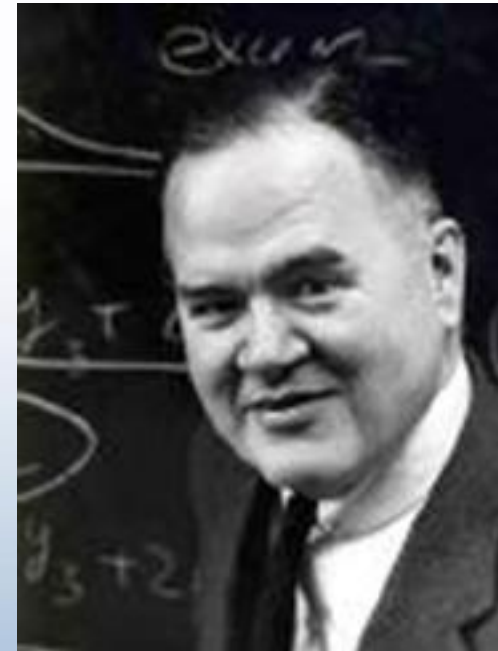
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BIT – BInary digiT

- *Bit* is a binary (base 2) digit: 0 or 1.
- Bits may be used to represent truth values.
- By convention:
0 represents “False”;
1 represents “True”.



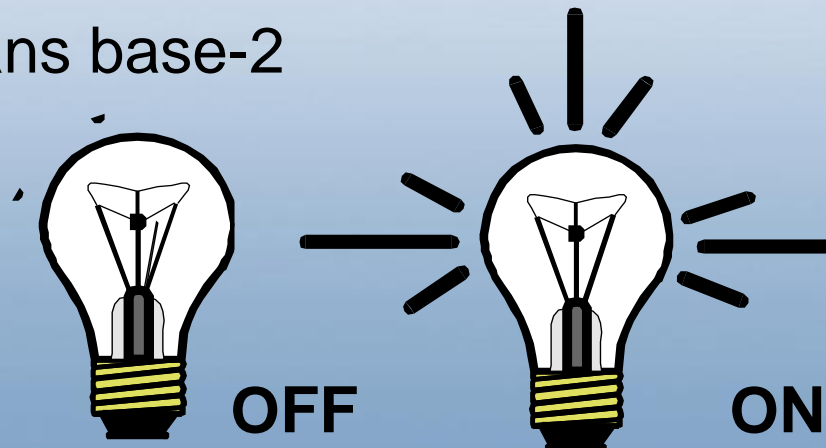
John Tukey
(1915-2000)

This terminology was introduced by statistician John Tukey in 1946.



BIT – BInary digiT

- Computers are made of a series of switches
- Each switch has two states: ON or OFF
- **Bit** (Binary Digit) = Basic unit of information, representing one of two discrete states.
The smallest unit of information within the computer.
- The only thing a computer understands.
- Bit has one of two values: 1 (ON) or 0 (OFF)
- Binary means base-2





Tables for Bit Operations

- **Boolean algebra** is like ordinary algebra except that variables stand for bits; + means “or”; multiplication means “and”.
- A variable is called **Boolean variable** if its value is either true or false.

NOT	0	1
	1	0

OR	0	1
0	0	1
1	1	1

AND	0	1
0	0	0
1	0	1

XOR	0	1
0	0	1
1	1	0



Bit Strings

- **Bit string of length n** is an ordered sequence (series, tuple) of $n \geq 0$ bits.
- By convention, bit strings are (sometimes) written left to right:
- the “first” bit of the bit string of length ten “1001101010” is 1.
- When a bit string represents a base-2 number, by convention, the first (leftmost) bit is the *most significant* bit.
- Example: $1101_2 = 8 + 4 + 1 = 13$.



Bitwise Operations

- Boolean operations can be extended to operate on bit strings as well as single bits.

❖ E.g.:

01 1011 0110

11 0001 1101

11 1011 1111 Bit-wise OR

01 0001 0100 Bit-wise AND

10 1010 1011 Bit-wise XOR



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Tautologies

Tautology is a compound proposition that is **true** no matter what the truth values of its atomic propositions are! When every row of the truth table gives T.

Example: $p \vee \neg p$

❖ What is its truth table?



Contradiction and Contingency

Contradiction is a compound proposition that is **false** no matter what!

When every row of the truth table gives F

Example: $p \wedge \neg p$

❖ Truth table?

A proposition that is neither a tautology nor a contradiction is called a **contingency**.



Logical Equivalence

- Compound proposition p is **logically equivalent** to compound proposition q , written $p \Leftrightarrow q$, if p and q contain the same truth values in all rows of their truth tables
- They express the same truth function (the same function **from** values for atoms **to** values for the whole formula):
“ p if and only if q ” or
“if p then q and conversely”

p	q	$p \Leftrightarrow q$
F	F	T
F	T	F
T	F	F
T	T	T



Proving Equivalence via Truth Tables

Example: prove that $p \vee q \Leftrightarrow \neg(\neg p \wedge \neg q)$.

p	q	$p \vee q$	$\neg p$	$\neg q$	$\neg p \wedge \neg q$	$\neg(\neg p \wedge \neg q)$
F	F	F	T	T	T	F
F	T	T	T	F	F	T
T	F	T	F	T	F	T
T	T	T	F	F	F	T



Equivalence Laws

- **Equivalence rules** provide a pattern or template that can be used to match all or part of a much more complicated proposition and to find an equivalence for it.
- These rules are similar to the arithmetic identities you may have learnt in algebra, but for propositional equivalences instead.



- *Identity:* $p \wedge \mathbf{T} \Leftrightarrow p$ $p \vee \mathbf{F} \Leftrightarrow p$
- *Domination:* $p \vee \mathbf{T} \Leftrightarrow \mathbf{T}$ $p \wedge \mathbf{F} \Leftrightarrow \mathbf{F}$
- *Idempotent:* $p \vee p \Leftrightarrow p$ $p \wedge p \Leftrightarrow p$
- *Double negation:* $\neg\neg p \Leftrightarrow p$
- *Commutative:* $p \vee q \Leftrightarrow q \vee p$ $p \wedge q \Leftrightarrow q \wedge p$
- *Associative:* $(p \vee q) \vee r \Leftrightarrow p \vee (q \vee r)$
 $(p \wedge q) \wedge r \Leftrightarrow p \wedge (q \wedge r)$

Equivalence Laws



Augustus De Morgan
(1806-1871)

- *Distributive:*

$$p \vee (q \wedge r) \Leftrightarrow (p \vee q) \wedge (p \vee r)$$

$$p \wedge (q \vee r) \Leftrightarrow (p \wedge q) \vee (p \wedge r)$$

- *De Morgan's laws:*

$$\neg(p \wedge q) \Leftrightarrow \neg p \vee \neg q$$

$$\neg(p \vee q) \Leftrightarrow \neg p \wedge \neg q$$



Derivation

Derivation is a finite sequence of propositions each of which follows from the preceding one in the sequence following equivalence laws:

$$p \Leftrightarrow q \Leftrightarrow r \Leftrightarrow \dots$$

Example:

$$(w \vee x) \vee (w \vee z) \Leftrightarrow \text{associative}$$

$$((w \vee x) \vee w) \vee z \Leftrightarrow \text{commutative}$$

$$((x \vee w) \vee w) \vee z \Leftrightarrow \text{associative}$$

$$((x \vee (w \vee w)) \vee z) \Leftrightarrow \text{idempotent}$$

$$(x \vee w) \vee z$$



Review: Propositional Logic

- Atomic propositions: p, q, r, \dots
- Boolean operators: $\neg \wedge \vee \oplus \dots$
- Compound propositions: $s ::= (p \wedge \neg q) \vee r$
- Equivalences: $p \wedge \neg q \Leftrightarrow \neg(p \rightarrow q)$
- Proving equivalences using:
 - Truth tables.
 - Symbolic derivations: $p \Leftrightarrow q \Leftrightarrow r \dots$



Questions?